Downlink Non-Orthogonal Multiple Access Systems With an Untrusted Relay

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Introduction

- NOMA techniques offer solutions to spectrum scarcity and congestion problems.
- Key feature: efficient utilization of available resources serving multiple users *simultaneously*.
- Using a relay node can generally boost up achievable rates in wireless communication systems.
- What if the relay is untrusted?
- Use physical layer security techniques to secure NOMA users' data from the relay, and still benefit from its presence.

System Model

• Two-user SISO Gaussian broadcast channel:



• An **untrusted** half-duplex relay assists the BS:



• BS uses superposition coding:

$$x = \sqrt{\alpha P} s_1 + \sqrt{\bar{\alpha} P} s_2$$

where $0 \le \alpha \le 1$ and $\bar{\alpha} \triangleq 1 - \alpha$.

• Treating the relay as an eavesdropper achieves the following secrecy rates [1, Theorem 5]:

$$r_{s,1} = \left[\log \left(1 + |h_1|^2 \alpha P \right) - \log \left(1 + |h_r|^2 \alpha P \right) \right]^+ r_{s,2} = \left[\log \left(1 + \frac{|h_2|^2 \bar{\alpha} P}{1 + |h_2|^2 \alpha P} \right) - \log \left(1 + \frac{|h_r|^2 \bar{\alpha} P}{1 + |h_r|^2 \alpha P} \right) \right]^+$$

• Can we achieve higher secrecy rates by employing the *untrusted* relay?

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Passive User Mode

- Communication occurs over two phases.
- Phase 1: BS broadcasts to users and relay.
- Phase 2: relay uses either *compress-and-forward* or *amplify-and-forward*.
- Both users *passively* listen to the communication.

Using *compress-and-forward*, the following secrecy rates are achievable with *passive users*: $r_{s,1}^{CF,P} = \frac{1}{2} \left| \log \left(1 + |h_1|^2 \alpha \bar{P} + \frac{|h_r|^2 \alpha \bar{P}}{1 + \sigma_Q^2} \right) - \log \left(1 + |h_r|^2 \alpha \bar{P} \right) \right|^2$ $r_{s,2}^{CF,P} = \frac{1}{2} \left[\log \left(1 + \frac{|h_2|^2 \bar{\alpha} \bar{P}}{1 + |h_2|^2 \alpha \bar{P}} + \frac{|h_r|^2 \bar{\alpha} \bar{P}}{1 + |h_r|^2 \alpha \bar{P} + \sigma_O^2} \right) - \log \left(1 + \frac{|h_r|^2 \bar{\alpha} \bar{P}}{1 + |h_r|^2 \alpha \bar{P}} \right) \right|^{\top}$

 $P \leq P$ is the new BS power, and σ_Q^2 is the quantization (compression) noise, whose value is such that decodability at *both* users is guaranteed (a function of $P - \overline{P}$). Active User Mode

- Phase 1: BS broadcasts to users and relay; users transmit a jamming signal to confuse the relay.
- Nodes are half-duplex \Rightarrow two-hop network.

Using <i>comp</i>	ess-and-forward, the following secrecy rates	1
	$r_{s,1}^{CF,A} = \frac{1}{2} \left[\log \left(1 + \frac{ h_r ^2 \alpha \bar{P}}{1 + \sigma_Q^2} \right) - \log \left(1 + \frac{1}{2} + \frac{1}$	1
	$r_{s,2}^{CF,A} = \frac{1}{2} \left[\log \left(1 + \frac{ h_r ^2 \bar{\alpha} \bar{P}}{1 + h_r ^2 \alpha \bar{P} + \sigma_Q^2} \right) - \right]$]

 $\delta \leq P - \overline{P}$ is the users' jamming power, and $\boldsymbol{g} \triangleq [g_1, g_2]$. Passive User Mode



Dashed lines are when relay is further than users from BS.









[1] E. Ekrem and S. Ulukus, "The secrecy capacity region of the Gaussian MIMO multi-receiver wiretap channel," *IEEE Trans. Inf. Theory*, vol. 57, no. 4, pp. 2083–2114, April 2011.

[2] X. He and A. Yener, "Cooperation with an untrusted relay: A secrecy perspective," IEEE Trans. Inf. Theory, vol. 56, no. 8, pp. 3807–3827, August 2010.

[3] A. Zewail and A. Yener, "Multi-terminal two-hop untrusted-relay networks with hierarchical security guarantees," IEEE Trans. Inf. Forensics Security, vol. 12, no. 9, pp. 2052–2066, September 2017.