# Lecture 11 Quantization

- Quantization
- Pulse-Code Modulation (PCM) and Line Codes
- Quantization Error
- Nonuniform Quantization and Companding

#### Contents

#### Quantization

- Pulse-Code Modulation (PCM) and Line Codes
- Quantization Error

• Nonuniform Quantization and Companding

# Analog to Digital (A/D) Conversion

- Analog to digital conversion steps
  - Step 1: sampling analog sources in discrete times
  - Step 2: quantizing the samples to discrete levels

**Quantization** is the process of transforming the sample amplitude of a baseband signal into a discrete amplitude taken from a finite set of possible levels.

• Quantization is an irreversible process.

- A sequence of samples is not a digital signal because the sample values can potentially take on a continuous range of values
- To complete analog to digital conversion, each sample value is mapped to a discrete level in a process called quantization
- Each discrete level is represented by a sequence of bits
- In a *b*-bit quantizer, each quantization level is represented with *b* bits, so that the number of levels equals  $L = 2^b$

# Example: Analog to digital (A/D) Conversion



• Here, we are mapping to the nearest limit (quantization level)

# **Uniform Quantizers**

In **uniform quantizers** representation levels (output levels) are uniformly spaced.



Two types of **uniform** quantization: (a) midtread (b) midrise



- Quantization
- Pulse-Code Modulation (PCM) and Line Codes
- Quantization Error

• Nonuniform Quantization and Companding

# Pulse-Code Modulation (PCM)

- PCM is the most basic form of digital pulse modulation
- In PCM, a message signal is represented by a sequence of coded pulses
- ${\scriptstyle \bullet} \,$  PCM usually uses only two pulse values, which represent 0 and 1



A PCM transmitter

In reality, PCM is *source coding* strategy, whereby an analog signal is converted into digital form. Once a sequence of 1s and 0s is produced, a *line code* is needed for electrical representation of that binary sequence.

#### Example:PCM process



The PCM process. [Taub & Schilling]

## Line Codes

A **line code** is used for an electrical representation of a binary data sequence (1s and 0s), say those of the PCM.

Examples: nonreturn-to-zero (NRZ), return-to-zero (RZ), etc.



# Line Codes

- (a) In **Unipolar NRZ** or on-off signaling symbol 1 is represented by transmitting a pulse of constant amplitude for the duration of the symbol  $(T_b)$ , and symbol 0 is represented by switching off the pulse
- (b) In **Polar NRZ** signaling symbols 1 and 0 are represented by transmitting pulses of amplitudes +A and -A.
- (c) In Unipolar RZ signaling symbol 1 is represented by transmitting a pulse of amplitude +A for half-symbol period, and symbol 0 is represented by transmitting no pulse.
- (d) Bipolar RZ has three amplitude levels. Specifically, pulses of amplitudes +A and -A are used alternately for symbol 1, and no pulse is transmitted for symbol 0. (it has a better error detection)
- (e) In Manchester Code symbol 1 is represented by a positive half-pulse of amplitude +A followed by a negative half-pulse of amplitude -A, with both pulses being half-symbol width. For symbol 0, the polarities of these two pulses are reversed.

#### **PCM Applications - Digital Telephone**

Audio signal bandwidth is about 15 kHz. But, for speech articulation 3.4 kHz is enough.<sup>1</sup>

- PCM transmission in digital telephone
  - 1. Use an LPF to eliminate components above 3.4 kHz
  - 2. Sample the signal at a rate of 8,000 sample/sec
  - 3. Quantize each sample to 256 levels ( $L = 256 \Rightarrow b = \log_2 L = 8$  bits/sample)



Thus, a telephone signal requires  $8 \times 8000 = 64,000$  binary pulses per second. That is, the bandwidth of the digital telephone is 64 kbps.

Note that the analog audio signal has a bandwidth of 3.4 kHz while its corresponding digital signal is 64 kilobits per second (kbps).

 $<sup>{}^1{\</sup>bf Q}{:}$  (Why) do you think people sound different on the phone from in person? Lecture 11: Quantization

# PCM Applications - Compact Disc (CD)

CD is a more recent application of PCM. CDs use a sampling rate of 44.1  $\rm kHz^1$  with 16-bit quantization for each sample.

**Example:** Storage capacity of CD

An audio CD holds up to 74 minutes, 33 seconds of sound. When the CD was first introduced in 1983, every 8 bits of digital signal data were encoded as 17 bits of signal and error correction data together. Given that 1 byte = 8 bits and that 1 megabyte  $(MB) = 2^{20}$  bytes<sup>2</sup>, show that the capacity of a compact disc is about 800 MB.

<sup>&</sup>lt;sup>1</sup>High-fidelity audio signal bandwidth is up to 20 kHz (20 kHz is the highest frequency generally audible by humans).

 $<sup>^{2}1</sup>$  MB is 1,024 kilobytes, or 1,048,576 (1024x1024) bytes, not one million bytes. Lecture 11: Quantization

## PCM Applications - Compact Disc (CD)

CD Capacity = 
$$(74 * 60 + 33)s \times 44, 100 \frac{\text{samples}}{s} \times 16 \frac{\text{bits}}{\text{samples}} \times \frac{17}{8} \times \frac{1}{8} \frac{\text{byte}}{\text{bits}} \times \frac{1}{2^{20}} \frac{\text{MB}}{\text{byte}} = 799.5 \text{MB}$$

 $^1\text{Originally}$  megabyte was used to describe a byte multiple (2  $^{20}$  = 1024  $\times$  1024) in computer programming.

1



• Quantization

- Pulse-Code Modulation (PCM) and Line Codes
- Quantization Error

• Nonuniform Quantization and Companding

#### **Uniform Quantizer's Error**



- Uniform quantization with L levels of a signal with peak amplitude  $m_p$  (i.e.,  $-m_p \le m(t) \le m_p$ ) has
  - ${\scriptstyle \circ }$  quantization step size is equal to  $\Delta$  =  $\frac{2m_p}{L}$
  - maximum quantization error =  $\frac{\Delta}{2} = \frac{m_p}{L}$
  - mean square error (or, quantization noise):  $N_q = \frac{\Delta^2}{12} = \frac{m_p^2}{3L^2}$

#### **Uniform Quantizer's SNR**

- Quantizer error is a random variable in the range:  $-\frac{\Delta}{2} \le q_e \le \frac{\Delta}{2}$
- We assume quantizer error is *uniformly* distributed, i.e., it is equally likely to lie in the range  $\left[-\frac{\Delta}{2} \ \frac{\Delta}{2}\right]$
- The variance (power) of error is given by<sup>1</sup>

$$\sigma_{q_e}^2 = \int_{-\frac{\Delta}{2}}^{\frac{\Delta}{2}} q_e^2 \frac{1}{\Delta} dq_e = \dots$$
$$= \frac{\Delta^2}{12}$$

• quantization step size is equal to  $\Delta = \frac{2m_p}{L}$ • mean square error:  $N_q = \frac{\Delta^2}{12} = \frac{m_p^2}{3L^2}$ 

<sup>1</sup>The variance of random variable x with pdf of f(x) is  $\sigma_x^2 = \int_{-\infty}^{\infty} x^2 f(x) dx$ . Lecture 11: Quantization

11 - 16

# **Uniform Quantizer's SNR**

• Thus, we get

Signal-to-noise ratio (SNR) for a uniform quantizer is obtained by

$$SNR = \frac{S}{N} = \frac{\text{signal power}}{\text{quantization noise power}} = 3L^2 \frac{\overline{m^2(t)}}{m_p^2}$$

where  $S = \overline{m^2(t)}$  is the power of the message signal m(t).

Note that signal power is the same as we defined earlier in the course

### **Uniform Quantizer's Error**

**Example:** signal quantized to 8 and 16 levels ( $L_1 = 8$  and  $L_2 = 16$ )



#### **Uniform Quantizer's Error**





#### **Uniform Quantization SNR**

**Example:** Consider the special case of a full-load sinusoidal modulating signal of amplitude  $A_m$  which utilizes all the representation levels provided. Find the SNR as a function rate in decibels.

• Signal power is  $S = \frac{A_m^2}{2}$ 

• 
$$m_p = A_m$$
,  $\Delta = \frac{2m_p}{L}$ , and  $L = 2^b$ 

• Quantization error (noise) power is  $\sigma_{q_e}^2 = \frac{\Delta^2}{12} = \frac{A_m^2}{3L^2}$ 

• Thus, SNR = 
$$\frac{S}{\sigma_{q_e}^2} = \frac{3}{2}L^2 = \frac{3}{2}2^{2b}$$

• SNR<sub>dB</sub> = 
$$10 \log_{10}$$
 SNR =  $1.76 + 6.02b \approx 1.8 + 6b$ 

In uniform quantizer, increasing one bit in the codeword quadruples the output SNR (6dB gain)

### **Blank Page**



• Quantization

- Pulse-Code Modulation (PCM) and Line Codes
- Quantization Error

• Nonuniform Quantization and Companding

# Why Nonuniform Quantization?





- Smaller amplitudes predominate in speech and larger amplitudes are much less frequent. This means the SNR will be low most of the time.
- Recall that SNR is an indication of the quality of the received signal.

## Why Nonuniform Quantization?

- The root of this difficulty (low SNR) lies in the fact that the quantizing steps are of uniform value  $(\Delta = \frac{2m_p}{L})$ .
- The quantization noise is directly proportional to the square of the step size  $(\sigma_{q_e}^2 = \frac{\Delta^2}{12})$ .
- The problem can be solved by using smaller steps for smaller amplitudes (more frequent amplitudes) such that quantization noise is smaller and thus SNR increases
- Such a quantizer will be nonuniform.

#### Nonuniform Quantizer

In **nonuniform quantizers** representation levels are spaced *variably*. Specifically, smaller steps are used for smaller amplitudes.



nonuniform quantization

#### **Nonuniform Quantizer**



nonuniform quantization

# Implementation of a Nonuniform Quantizer using a Uniform Quantizer

- A nonuniform quantizer can be realized by cascading a signal compressor and uniform quantizer.
- Logarithmic curve (right figure) is a compressor



(a) nonuniform quantization (varying quantization levels)
(b) implementation of nonuniform quantizer by cascading a signal compressor (logarithmic curve) and uniform quantizer

# **Nonuniform Quantizers**



- A nonuniform quantizer can be realized by cascading a signal compressor and uniform quantizer.
- To restore the signal samples to their correct relative level, we need to compensate for the compression. This device is called expander.
- Expander is done at the receiver side.
- Compressor + expander = compander

## **Companding Curves**



#### $\mu$ -Law & A-Law Companding

- Logarithmic compression can be approximated by  $\mu$ -law and A-law
- Let  $-m_p \le m(t) \le m_p$
- Define x as normalized input, i.e.,  $x = \frac{m}{m_n}$
- y is the normalized output, i.e., actual output divided by  $m_p$
- 1.  $\mu$ -law: (North America and Japan)

$$y = \operatorname{sign}(x) \frac{\ln(1+\mu|x|)}{\ln(1+\mu)} \qquad 0 \le |x| \le 1$$

•  $\mu = 255$  (for 8-bit codes) is used in digital telephone in NA

 when µ → 0, using L'Hospital's rule, you can check that y = sign(x)|x| = x, which implies a uniform quantizer

#### $\mu\text{-Law}$ & A-Law Companding

2. A-law: (Europe & rest of the world)

$$y = \begin{cases} \operatorname{sign}(x) \frac{A|x|}{1+\ln A} & 0 \le |x| \le \frac{1}{A} \\ \\ \operatorname{sign}(x) \frac{1+\ln A|x|}{1+\ln A} & \frac{1}{A} \le |x| \le 1 \end{cases}$$

- The standard value is A = 87.7
- · A-law is used for international connections if at least one country uses it

#### Comparison of $\mu$ -Law and A-Law



Lecture 11: Quantization

#### **Example:** $\mu$ -Law Companding

**Example:** In a compander with a maximum voltage range of 2V and  $\mu$  = 255, what are the output voltage and gain for

- i. the input voltage of 0.5 V
- ii. the input voltage of 0.1  ${\sf V}$

i. 
$$V_{\text{in}} = 0.5 \implies x = \frac{V_{\text{in}}}{V_p} = \frac{0.5}{2} = 0.25 \implies y = \frac{\ln(1+\mu|x|)}{\ln(1+\mu)} = 0.75$$
  
 $y = \frac{V_{\text{out}}}{V_p} \implies V_{\text{out}} = 0.75V_p = 1.5V \implies \text{Gain} = \frac{y}{x} = 3$ 

ii.  $V_{\text{in}} = 0.1V \implies x = 0.05 \implies y = 0.47 \implies \text{Gain} = \frac{0.47}{0.05} = 9.4$ 

• As expected, there is a much higher gain for lower inputs

#### The Effect of $\mu$ -Law on Signal Histogram



• the above is the histogram of an audio signal (we will see in the lab)

• companding decreases the number of low-amplitude inputs

#### $\mu\text{-Law's}$ SNR



- $\mu$ -Law's SNR happens to be  $\left| \frac{S_o}{N_o} \approx \frac{3L^2}{[\ln(1+\mu)]^2} \right|$  for  $\mu^2 \gg \frac{m_p^2}{\overline{m^2(t)}}$
- The output SNR for  $\mu = 255$  and  $\mu = 0$  (uniform quantization) as a function of the message signal power is shown above

### **Example:** $\mu$ -Law Companding

**Example:** Consider a speech signal with spectral components in the range of 300 to 3000 Hz. Recall from lecture notes 5 that speech has a high crest factor as shown in the figure.



Assume that a sampling rate of 8,000 samples per second will be used to generate a PCM signal. The required SNR is 30 dB.

- i. what is the minimum number of bits per sample needed.
- ii. repeat part i when a  $\mu\text{-}\mathsf{law}$  compander with  $\mu$  = 255 is used.
- iii. which one is more efficient?

# Solution

i. uniform quantizer:

Let  $L = 2^b$  be the number of quantization levels. For  $x_p \simeq 35 x_{\rm rms}$  (see the figure), and a uniform quantizer we have

$$SNR_o = \frac{\text{signal power}}{\text{quantization noise power}} = \frac{x_{\text{rms}}^2}{\frac{\Delta^2}{12}} = \frac{\left(\frac{x_p}{35}\right)^2}{\left(\frac{2x_p}{L}\right)^2}}{\frac{\left(\frac{2x_p}{L}\right)^2}{12}} = \frac{3}{35^2}L^2.$$

Then, SNR in dB is given by

$$(SNR_o)_{dB} = 10 \log_{10} SNR_o = 6b - 26.1$$

Next, 6b - 26.1 > 30 results in b > 9.35. Thus, the smallest number of bits is b = 10.

## Solution

ii.  $\mu$ -Law companding (non-uniform quantizer):

For  $\mu^2 \gg \frac{m_p^2}{\overline{m^2(t)}}$  , we know that  $\mu\text{-Law's SNR}$  is given by

$$\mathsf{SNR}_o = \frac{3L^2}{[\ln(1+\mu)]^2}$$

Then,

 $(\mathsf{SNR}_o)_{\mathrm{dB}} = 10 \log_{10} \mathsf{SNR}_o = 6b + 10 \log_{10} \frac{3}{[\ln(1+\mu)]^2} = 6b - 10.1$ 

Next, 6b - 10.1 > 30 results in b > 6.68; i.e., b = 7. Note that the condition is satisfied here  $(255^2 \gg 35^2)$ .

iii. clearly non-uniform quantization is more efficient as it needs fewer bits to represent each quantized sample.