Lecture 12 Performance Analysis of Digital Systems

- Noise in Communication Systems
- Optimum Detection for Binary Signals
- Probability of Error
- Summary
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Noise

Noise is an unwanted wave that disturbs the transmission of signals.

Where does noise come from?

1. External sources: e.g.,

- Industrial (electric motors, generators, automotive ignition systems, etc.)
- · Atmospheric, often referred to as static, usually comes from lightning
- *Extraterrestrial* noise (solar and cosmic) comes from sources in space, e.g., sun and other stars)

-The sun produces an awesome amount of noise that, at its peak, makes some frequencies unusable for communication.

- 2. **Internal sources:** generated by communication devices themselves. This type of noise represents a basic limitation on the performance of electronic communication systems.
 - **Thermal noise:** caused by a phenomenon known as thermal agitation, the random motion of free electrons in a conductor caused by heat.
 - **Shot noise:** the electrons are discrete and are not moving in a continuous steady flow, so the current is randomly fluctuating.

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Noise

- When you turn on an AM or FM receiver and tune it to some position between stations, the hiss or static that you hear in the speaker is noise.
- The noise level in a system is proportional to temperature, bandwidth, the amount of current flowing in a component, the gain of the circuit, and the resistance of the circuit.
- Noise can be external to the receiver or originate within the receiver itself.
- Noise is a problem in communication systems whenever the received signals are very low in amplitude.
- Regardless of its source, noise shows up as a random ac voltage and can be seen on an oscilloscope.

Noise that occurs in transmitting digital data causes bit errors.

Thermal Noise

• The mean-square value of the thermal noise voltage measured across the terminals of the resistor of *R* ohms equals

 $\mathbb{E}[V^2] = 4kTRB_N \quad \text{volts}^2$

- k is Boltzmann's constant $1.38 \times 10^{-23} watts/Hz/K$
- ${\, \bullet \, } T$ is absolute temperature in degrees Kelvin
- B_N is the measured bandwidth in Hertz
- What if we have multiple resistors?



Modeling a noisy resistor with a noiseless resistor pulse noise source

Thermal Noise - Example

Example: Noise Voltage in AM Radio

The front-end filter of a radio passes the broadcast AM band from 535 kHz to 1605 kHz. The radio input has an effective resistance of 300 ohms. What is the root-mean-square (rms) noise voltage that we would expect to observe due to this resistance?

- The bandwidth is 1605 535 = 1070 kHz.
- The mean-square thermal noise voltage at room temperature (290 Kelvin) is

$$\mathbb{E}[V^2] = 4kTRB_N = 4 \times 1.38 \times 10^{-23} \times 290 \times 300 \times 1070 \times 10^3$$
$$= 5.14 \times 10^{-12} \text{ V}^2$$

• Hence, the rms thermal noise voltage is $\sqrt{5.14 \times 10^{-12}}$ = 2.3 microvolts.

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Thermal Noise Power Values - Examples

Thermal noise power across a resistor $P = \mathbb{E}[V^2]/R$ (source: Wikipedia)

Bandwidth (Δf)	Thermal noise power at 300 K (dBm)	Notes
1 Hz	-174	
10 Hz	-164	
100 Hz	-154	
1 kHz	-144	
10 kHz	-134	FM channel of 2-way radio
100 kHz	-124	
180 kHz	-121.45	One LTE resource block
200 kHz	-121	GSM channel
1 MHz	-114	Bluetooth channel
2 MHz	-111	Commercial GPS channel
3.84 MHz	-108	UMTS channel
6 MHz	-106	Analog television channel
20 MHz	-101	WLAN 802.11 channel
40 MHz	-98	WLAN 802.11n 40 MHz channel
80 MHz	-95	WLAN 802.11ac 80 MHz channel
160 MHz	-92	WLAN 802.11ac 160 MHz channel
1 GHz	-84	UWB channel

Shot Noise (or Poisson Noise)

- Current flow in any device is not direct and linear.
- The current carriers, electrons or holes, sometimes take random paths from source to destination (e.g., collector or drain in a transistor)
- It is this random movement that produces the shot effect.
- The amount of shot noise is directly proportional to the amount of dc bias flowing in a device.
- The rms of noise current in a device I_n is given by $I_n = \sqrt{2qIB_N} A$
 - q is charge on an electron $(1.6 \times 10^{-19} C)$
 - I is direct current, A
 - ${\scriptstyle \bullet \ } B_N$ is the measured bandwidth, Hz

Shot Noise - Example

Example: What is noise current for a dc bias of 0.1 mA and a bandwidth of 12.5 kHz.

$$\begin{split} I_n = \sqrt{2qIB_N} &= \sqrt{2\times 1.6\times 10^{-19}\times 0.1\times 10^{-3}\times 12.5\times 10^3} \\ &= 0.623\times 10^{-9} \text{ A} \end{split}$$

Hence, the noise current is 0.623 nA.

White Noise

The noise analysis of communication systems is often based on an *idealized* noise process called white noise.

- The *power spectral density* (PSD) of white noise is constant in frequency.
- White noise is analogous to the term white light in the sense that all frequency components are present in.
- PSD of white noise w(t) will be denoted by $S_W(f) = \frac{N_0}{2}$



White Noise Examples

- Thermal noise and shot noise are examples of white noise processes generated by electrical circuits.
- The amplitude of the noise voltage is unpredictable.
- Both have a zero-mean Gaussian distribution (following from the central limit theorem)



Power Spectral Density, $S_{u}(f)$

White Noise Examples - Thermal Noise

Thermal noise is white, and has the following characteristics:

- Fluctuations are Gaussian distributed in time.
- Fluctuations in time are uncorrelated.
- Flat bandwidth. The Fourier transform of thermal noise is a straight line (DC).



Signal fluctuations due to thermal noise in the time domain

Filtering the White Noise

What happens if we filter a white noise?

- If a noise signal is fed into a selective tuned circuit, many of the noise frequencies are rejected and the overall noise level goes down.
- Assume the white noise w(t) (whose PSD is $\frac{N_0}{2}$) is filter by a filter whose transfer function is H(f).
- Let's call the filtered noise n(t). The PSD of the filtered noise will be

$$S_N(f) = \frac{N_0}{2} |H(f)|^2$$

• The average power of the filtered noise is given by

$$\mathbb{E}[n^2(t)] = \frac{N_0}{2} \int_{-\infty}^{\infty} |H(f)|^2 df$$

Gaussian Random Variable

The probability density function (pdf) of a Gaussian random variable Z with a mean μ and a variance σ^2 is defined as

$$f_Z(z) = \frac{1}{\sqrt{2\pi\sigma^2}} e^{-\frac{(z-\mu)^2}{2\sigma^2}}, \qquad -\infty < z < \infty$$

• for $\mu = 0$ and $\sigma^2 = 1$ we get $f_Z(z) = \frac{1}{\sqrt{2\pi}}e^{-\frac{z^2}{2}}$ which is called *standard* Gaussian pdf



Q function

• For a given pdf $f_Z(z)$, by definition,

$$P(a < Z \le b) = \int_a^b f_Z(z) dz$$

 $P(a < Z \leq b)$ is the probability that Z is less than b and greater than a

- For a Gaussian pdf, such an integral cannot be evaluated in a closed form and must be evaluated numerically
- Numerical tables usually are available for *Q* function which is defined as

$$Q(x) = \int_x^\infty \frac{1}{\sqrt{2\pi}} e^{-\frac{z^2}{2}} dz$$



Q function

- Clearly, Q(x) is a decreasing function (as x increases, Q(x) decreases)
- ${\ensuremath{\, \bullet }}$ For example, using Q-function tables or ${\rm MATLAB},$ we get
 - Q(-1.22) = 0.8888
 - Q(2.5) = 0.0062
 - Q(3) = 0.0013
 - $Q(5) = 2.87 \times 10^{-7}$
- In MATLAB, you can use the command
 - qfunc(x) to find the Q function of x, and
 - *qfuncinv(x)* to find its inverse
- Some special values
 - $Q(\infty) = \ldots$
 - $Q(-\infty) = \dots$
 - Q(0) = ...
 - probability that x falls outside [-3, 3] is 2Q(3) = 0.0026

Gaussian Noise and AWGN

Gaussian noise is a random process that has a Gaussian distribution in the time domain. That is, the probability that specific amplitudes of noise will occur follows a Gaussian distribution curve.

Property 1: If a Gaussian noise is applied to a stable linear filter, then the random process developed at the output of filter will also be Gaussian.

In communications, noise is typically assumed to be additive white Gaussian noise (AWGN).

- Additive: noise is added to the signal
- White: noise has uniform power across the frequency band
- Gaussian: noise has a Gaussian distribution in the time domain

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Performance Metrics in Analog and Digital Communications

- Analog communications
 - Objective: achieving high fidelity in waveform reproduction
 - Performance criterion: output signal-to-noise ratio (SNR)
- Digital communications
 - **Objective:** accurately determine the *transmitted symbol* (NOT to reproduce the waveform that carries the symbol with fidelity)
 - **Performance criterion:** probability of symbol (or bit) error; also known as *bit error rate (BER)*

Main Questions (this lecture):

- What is the structure of *optimum detection receivers* to minimize BER?
- How to analyze BER performance?

Any Suggestions for Receiver Structure?



(a) Noise on a binary signal. (b) **Example** of transmitted binary signal and received one at SNR=0 dB.

Q: How can we guess if '0' is transmitted or '1'? Any suggestions about receiver structure?

Receiver Model



- Received signal (filter input) x(t) = g(t) + w(t) $0 \le t \le T$
 - g(t) is a general pulse representing 0 or 1
 - $\bullet \ T$ is the duration of the pulses
 - w(t) is a sample function of a *white noise* with zero mean and PSD $rac{N_0}{2}$
- The function of receiver is to detect g(t) from the received signal x(t)

When no pulse is transmitted for '0', the purpose of detection is to establish the presence or absence of an information-bearing signal in noise.

Received SNR

Since the filter is linear, the output signal y(t) can be expressed as

$$x(t) = g(t) + w(t)$$
Linear time-
invariant filter of
impulse response
 $h(t)$
 $y(t) = g_o(t) + n(t)$

• $g_o(t)$ and n(t) are the filtered signal and noise components • $g_o(t) = h(t) \star g(t)$ or equivalently $G_o(f) = H(f)G(f)$ • $g_o(t) = \mathcal{F}^{-1}[G_o(f)] = \int_{-\infty}^{\infty} H(f)G(f)e^{j2\pi ft}df$ • $\mathbb{E}[n^2(t)] = \frac{N_0}{2}\int_{-\infty}^{\infty} |H(f)|^2 df$ (average power of output noise)

- peak pulse SNR is defined as $SNR = \frac{|g_o(t)|^2}{\mathbb{E}[n^2(t)]}$ where $|g_o(t)|^2$ is the instantaneous power of the output signal
- when the filter output is sampled at t = T we have

$$\operatorname{SNR} = \frac{|g_o(T)|^2}{\mathbb{E}[n^2(t)]} = \frac{\left|\int_{-\infty}^{\infty} H(f)G(f)e^{j2\pi fT}df\right|^2}{\frac{N_0}{2}\int_{-\infty}^{\infty}|H(f)|^2df}$$

Received SNR- Critical Questions

Q1: Why filter's output is sampled at t = T not a different time t?

Q2: What H(f) maximizes the SNR?

Matched Filter

Using Cauchy-Schwarz inequality we can prove that ¹

$$\operatorname{SNR}_{\max} = \frac{2}{N_0} \int_{-\infty}^{\infty} |G(f)|^2 df$$

and, it is achieved by the following H(f):

$$H_{\mathsf{opt}}(f) = G^*(f)e^{-j2\pi fT}$$

• The time domain representation of the filter

- in general, $h_{\text{opt}}(t) = \mathcal{F}^{-1}[H_{\text{opt}}(f)] = \int_{-\infty}^{\infty} G^*(f) e^{-j2\pi f(T-t)} df$
- for a real signal g(t), we have $G^*(f) = G(-f)$ and we obtain²

$$h_{\mathsf{opt}}(t) = g(T - t)$$

The impulse response of **optimum filter** is time-reversed and delayed version of the input pulse g(t).

• This filter is called *matched filter*.

²In general, $h_{opt}(t) = k g(T - t)$ where k is a scaling factor.

¹See the Appendix for the proof.

Properties of Matched Filter

Recall that *matched filter* is a linear filter designed to provide maximum output SNR for a given transmitted signal (pulse).

- **Property 1:** $h_{opt}(t) = g(T t)$
- **Property 2:** The output SNR of a matched filter depends only on the ratio of the signal energy to the white noise psd at the filter input, i.e.,

$$SNR_{max} = \frac{E}{\frac{N_0}{2}} = \frac{2E}{N_0}$$

where $E = \int_{-\infty}^{\infty} |G(f)|^2 df$ is the energy of pulse signal g(t).

• Note that, with the above SNR_{max} , dependence on the waveform of the input g(t) is completely removed by the matched filter. That is,

Only energy of pulse g(t) is important, not its shape.

Matched Filter - Example 1



To design the MF, we only need to know the shape of the binary pulse.

Matched Filter - Example 2

- Transmitted pulse (Fig. a) Q1: What is the match filter for this pulse?
- MF output (Fig. b)
 Q2: Where does the maximum happens?

Q3: What is the duration of the output? Why so?

• Output of integrate-and-dump circuit (Fig. c)



Matched Filter-Special Case

• For the special case of rectangular pulse, the matched filter can be implemented using integrate-and-dump circuit



Line Codes and their Power Spectrum



Q: What are the MFs corresponding to each of these line codes?

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Binary Digital Signal Detection

Binary digital signal detection in short:



So far, we have talked about

- AWGN Channel
- Matched Filter (MF)

Example: Binary Detection with and without MF



Binary Detection



Two steps are involved in signal detection

- Step 1: reducing the received signal x(t) into a number y(T)
 - 1. find a linear filter that *maximizes* output SNR [Answer: *matched filter*]
 - 2. sample the output of filter at certain time (t = T) to get a number y(T)
- Step 2: comparing y(T) with a certain threshold and deciding which signal (i.e., $s_1(t)$ or $s_2(t)$) has been transmitted

We are going to discuss the last block (Step 2), in the following.

• Consider binary transmission with

$$s_i(t) = \begin{cases} s_1(t), & 0 \le t \le T, & \text{for } 1\\ s_2(t), & 0 \le t \le T, & \text{for } 0 \end{cases}$$

• The received signal x(t) can be represented as

$$x(t) = s_i(t) + w(t), \quad 0 \le t \le T, \quad i = 1, 2$$

where w(t) is a zero mean AWGN.

• After filtering the signal x(t) we can write

$$y(t) = a_i(t) + n(t), \quad 0 \le t \le T, \quad i = 1, 2$$

• After sampling at t = T, we get

$$y(T) = a_i(T) + n(T), \quad i = 1, 2$$

For notational convenience, we write this as

$$y = a_i + n, \quad i = 1, 2$$

Note that y is a Gaussian r.v. with mean a_1 (a_2) when $s_1(s_2)$ was sent.



- At this point, by comparing y with a threshold λ , we can decide whether 0 or 1 was sent. Specifically, $y \underset{s_2}{\overset{s_1}{\gtrless}} \lambda$.
- When $p(s_1) = p(s_2) = \frac{1}{2}$, the optimum threshold is $\lambda = \frac{a_1+a_2}{2}$. That is, the decision is made based on the distance of y from a_1 and a_2 . If y is closer to a_1 , we decide 1 (s_1) has been sent. Otherwise, we decide 0 (s_2) has been sent.
- Now, we can find the decision error; there are two possible error events
 1 (s₁) is chosen when 0 (s₂) was actually sent ≜ p_{e2}
 0 (s₂) is chosen when 1 (s₁) was actually sent ≜ p_{e1}
 Mathematically,

$$p_{e2} = \int_{\lambda}^{+\infty} f(y|s_2) dy$$
$$p_{e1} = \int_{-\infty}^{\lambda} f(y|s_1) dy$$



 a_2

• Finally, $P_e = p(s_1)p_{e1} + p(s_2)p_{e2}$

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• Clearly, $p_{e1} = p_{e2}$. Thus, when $p(s_1) = p(s_2) = \frac{1}{2}$ we get $P_e = p_{e1} = p_{e2}$

$$P_e = p_{e2} = \int_{\lambda}^{+\infty} f(y|s_2) dy$$
$$= \int_{\lambda}^{+\infty} \frac{1}{\sqrt{2\pi\sigma_n}} e^{-\frac{(y-a_2)^2}{2\sigma_n^2}} dy$$
$$= Q\left(\frac{a_1 - a_2}{2\sigma_n}\right)$$

where the last step is obtained by changing variable $x = \frac{y-a_2}{\sigma_n}$ and using the definition of Q-function.

• Now, let us define $\beta = \frac{a_1 - a_2}{\sigma_n}$. Then,

$$\beta^{2} = \frac{|a_{1} - a_{2}|^{2}}{\sigma_{n}^{2}} = \frac{|a_{1}(T) - a_{2}(T)|^{2}}{\mathbb{E}[n^{2}(t)]}$$

can be seen as the output SNR of the MF for an input $s_1(t) - s_2(t)$.

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• Thus,
$$\beta_{\max}^2 = \frac{2E_d}{N_0}$$
 where $E_d = \int_0^T [s_1(t) - s_2(t)]^2 dt$
• Finally, $P_e = Q(\frac{a_1 - a_2}{2\sigma_n}) = Q(\frac{\beta_{\max}}{2}) = Q(\sqrt{\frac{E_d}{2N_0}})$

Summary! $P_e = Q(\sqrt{\frac{E_d}{2N_0}})$ where E_d is the energy of difference of the two pulses; i.e., $E_d = \int_0^T [s_1(t) - s_2(t)]^2 dt$

 It is also very common to write P_e in terms of average energy per bit (E_b)

 E_b ≜ E₁+E₂/2 = ½[∫₀^T s₁²(t)dt + ∫₀^T s₂²(t)dt]

 Then, we just need to find the relationship

between E_b and E_d



Summary (Receiver Structure)



- In this structure, matched filter¹ is matched to $s_1(t) s_2(t)$
- Often, one of the pulses $(s_2(t))$ is 0 ۲

¹Equivalently, we can have two branches one matching to $s_1(t)$ and the other one matching to $s_2(t)$ Lecture 12: Performance Analysis

Probability of Error for Polar NRZ Signaling

Example: polar NRZ signaling
•
$$s_i(t) = \begin{cases} s_1(t) = A, & 0 \le t \le T, & \text{for } 1 \\ s_2(t) = -A, & 0 \le t \le T, & \text{for } 0 \end{cases}$$

• $E_d = \int_0^T [s_1(t) - s_2(t)]^2 dt = \int_0^T [2A]^2 dt = 4A^2T$
• $P_e = Q(\sqrt{\frac{E_d}{2N_0}}) = Q(\sqrt{\frac{2A^2T}{N_0}})$
• Since $E_b = \frac{1}{2} [\int_0^T s_1^2(t) dt + \int_0^T s_2^2(t) dt] = A^2T$, we can also write $P_e = Q(\sqrt{\frac{E_d}{2N_0}}) = Q(\sqrt{\frac{2A^2T}{N_0}}) = Q(\sqrt{\frac{2E_b}{N_0}})$

Example: For A = 5 mV, bit rate 10^3 bits/s (i.e., $T = 10^{-3}$), and noise power spectral density $\frac{N_0}{2} = 10^{-9}$ Watts/Hz, we get

$$P_e = Q\left(\sqrt{\frac{2A^2T}{N_0}}\right) = Q\left(\sqrt{\frac{2 \times 0.005^2 \times 10^{-3}}{2 \times 10^{-9}}}\right) = Q(5)$$

Using Q-function table, we get $P_e = Q(5) = 2.8665 \times 10^{-7}$.

Probability of Error for Unipolar NRZ Signaling

Example: unipolar NRZ signaling
•
$$s_i(t) = \begin{cases} s_1(t) = A, & 0 \le t \le T, \text{ for 1} \\ s_2(t) = 0, & 0 \le t \le T, \text{ for 0} \end{cases}$$

• $E_d = \int_0^T [s_1(t) - s_2(t)]^2 dt = \int_0^T A^2 dt = A^2 T = 2E_b \text{ (why?)}$
• $P_e = Q(\sqrt{\frac{E_d}{2N_0}}) = Q(\sqrt{\frac{A^2T}{2N_0}}) = Q(\sqrt{\frac{E_b}{N_0}})$
• Note that $E_1 = \int_0^T s_1^2(t) dt = A^2 T$ and $E_2 = 0$; thus, $E_b = \frac{A^2T}{2}$
• Assuming that $A = 5 \text{ mV}$, bit rate is 10³ bits/s, and noise power spectral density is $\frac{N_0}{2} = 10^{-9} \text{ Watts/Hz}$, we get
 $P_e = Q(\sqrt{\frac{A^2T}{2N_0}}) = Q(\sqrt{\frac{0.005^2 \times 10^{-3}}{4 \times 10^{-9}}}) = Q(2.5) = 6.2 \times 10^{-3}$

Q: Which signaling (unipolar NRZ or polar NRZ) is better in terms of probability of error and why?

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Bandpass Signaling vs. Baseband Signaling

- Thus far only baseband signaling has been considered: an information source is usually a baseband signal.
- Some communication channels have a bandpass characteristic, and will not propagate baseband signals.
- In these cases, modulation is required to impart the source information onto a bandpass signal with a carrier frequency f_c by the introduction of amplitude and/or phase perturbations
- The most common binary bandpass signaling techniques are
 - Amplitude-Shift Keying (ASK), a.k.a. On-Off Keying (OOK)
 - Phase-Shift Keying (PSK), a.k.a. binary PSK (BPSK)
 - Frequency-Shift Keying (FSK)

Bandpass Signaling







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Probability of Error for ASK in terms of Energy per Bit

It is more common to show P_e in terms of average energy per bit (E_b) rather than energy of the pulses difference (E_d)

Example: ASK signaling
•
$$s_i(t) = \begin{cases} s_1(t) = A \cos \omega_c t, & 0 \le t \le T, \text{ for } 1\\ s_2(t) = 0, & 0 \le t \le T, \text{ for } 0 \end{cases}$$
 and $\omega_c = n \frac{2\pi}{T}$.
• $E_d = \int_0^T [s_1(t) - s_2(t)]^2 dt = \int_0^T A^2 \cos^2 \omega_c t \, dt = \frac{A^2 T}{2}$
• $E_1 = \int_0^T [s_1(t)]^2 dt = \int_0^T A^2 \cos^2 \omega_c t \, dt = \frac{A^2 T}{2}$
• $E_2 = \int_0^T [s_2(t)]^2 dt = \int_0^T 0 \, dt = 0$
• $E_b = \frac{E_1 + E_2}{2} = \frac{A^2 T}{2} + \frac{A^2 T}{4} = \frac{E_d}{2}$
• $P_e = Q(\sqrt{\frac{E_d}{2N_0}}) = Q(\sqrt{\frac{E_b}{N_0}})$

Probability of Error for BPSK

Example: BPSK signaling
•
$$s_i(t) = \begin{cases} s_1(t) = A \cos \omega_c t, & 0 \le t \le T, & \text{for } 1 \\ s_2(t) = -A \cos \omega_c t, & 0 \le t \le T, & \text{for } 0 \\ \text{and } \omega_c = \frac{2\pi}{T}. \end{cases}$$

• $E_d = \int_0^T [s_1(t) - s_2(t)]^2 dt = \int_0^T 4A^2 \cos^2 \omega_c t dt = 2A^2T$
• $E_1 = \int_0^T [s_1(t)]^2 dt = \int_0^T A^2 \cos^2 \omega_c t dt = \frac{A^2T}{2}$
• $E_2 = E_1$
• $E_b = \frac{E_1 + E_2}{2} = \frac{A^2T}{2} = \frac{E_d}{4}$
• $P_e = Q(\sqrt{\frac{E_d}{2N_0}}) = Q(\sqrt{\frac{A^2T}{N_0}}) = Q(\sqrt{\frac{2E_b}{N_0}})$

• For the same E_b , PSK has a lower error rate than ASK. Why?

FSK Modulation/Signaling



Probability of Error for FSK

$$\begin{aligned} & \text{Example: FSK signaling} \\ & \quad s_i(t) = \begin{cases} s_1(t) = A\cos\omega_1 t, \quad 0 \leq t \leq T, \quad \text{for } 1\\ s_2(t) = A\cos\omega_2 t, \quad 0 \leq t \leq T, \quad \text{for } 0 \end{aligned} \\ & \quad E_d = \int_0^T [s_1(t) - s_2(t)]^2 dt = \int_0^T A^2 [\cos\omega_1 t - \cos\omega_2 t]^2 dt = \\ & A^2 \int_0^T [\cos^2\omega_1 t + \cos^2\omega_2 t - 2\cos\omega_1 t\cos\omega_2 t] dt = \\ & A^2 \int_0^T [\frac{1 + \cos 2\omega_1 t}{4\omega_1 T} + \frac{1 + \cos 2\omega_2 t}{4\omega_2 T} - \cos 2(\omega_1 - \omega_2) t - \cos 2(\omega_1 + \omega_2) t] dt = \\ & A^2 T [1 + \frac{\sin 2\omega_1 T}{4\omega_1 T} + \frac{\sin 2\omega_2 T}{4\omega_2 T} - \frac{\sin 2(\omega_1 - \omega_1) T}{2(\omega_1 - \omega_2) T} - \frac{\sin 2(\omega_1 + \omega_2) T}{2(\omega_1 + \omega_2) T}] \end{aligned} \\ & \quad \text{If } \omega_1 T \gg 1, \ \omega_2 T \gg 1, \ \text{and } |\omega_1 - \omega_2| T \gg 1, \ \text{then } E_d \approx A^2 T \ \text{and} \\ & P_e \text{ is approximated as} \end{aligned} \\ & \quad P_e = Q \Big(\sqrt{\frac{E_d}{2N_0}} \Big) \approx Q \Big(\sqrt{\frac{A^2 T}{2N_0}} \Big) = Q \Big(\sqrt{\frac{E_b}{N_0}} \Big) \end{aligned}$$
 \\ & \quad \mathbb{E}_b = \frac{E_1 + E_2}{2} \approx \frac{A^2 T}{2} = \frac{E_d}{2} \end{aligned}
$$& \quad \text{if } \omega_1 = n_1 \frac{\pi}{T} \ \text{and } \omega_2 = n_2 \frac{\pi}{T} \ \text{the 's' becomes '='. (why?)} \end{aligned}$$

Example: OOK

Example: OOK (ASK) signaling

An on-off binary system uses the pulse waveform

$$s_i(t) = \begin{cases} s_1(t) = A \sin \frac{\pi t}{T}, & 0 \le t \le T, & \text{for 1} \\ s_2(t) = 0, & 0 \le t \le T, & \text{for 0} \\ \text{with } A = 0.2 \text{ mv and } T = 2\mu s. \text{ For noise power spectral density} \\ \frac{N_0}{2} = 10^{-15} \text{ Watts/Hz} \end{cases}$$

- 1. Determine the probability of error.
- 2. What is the shape of the optimum detection filter?

1.
$$E_d = \int_0^T [s_1(t) - s_2(t)]^2 dt = \frac{A^2 T}{2}$$

 $P_e = Q(\sqrt{\frac{E_d}{2N_0}}) = Q(\sqrt{\frac{A^2 T}{4N_0}}) = Q(\sqrt{\frac{0.0002^2 \times 2 \times 10^{-6}}{4 \times 2 \times 10^{-15}}}) = 7.83 \times 10^{-4}$
2. $h_{\text{opt}}(t) = s_1(T - t) = A \sin \frac{\pi(T - t)}{T} = A \sin \frac{\pi t}{T}$

Probability of Error with Binary Signaling

- For any binary signaling with $p(0) = p(1) = \frac{1}{2}$, we have $P_e = Q(\sqrt{\frac{E_d}{2N_0}})$
- Typically, we are interested in P_e in terms of E_b , average energy per bit. This is shown for various baseband and passband binary signaling below.

Signaling Type	Signaling Name	P_e
baseband	polar NRZ	$Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$
baseband	unipolar NRZ	$Q(\sqrt{\frac{E_b}{N_0}})$
baseband	polar RZ	$Q\left(\sqrt{\frac{E_b}{N_0}}\right)$
baseband	unipolar RZ	$Q(\sqrt{\frac{E_b}{N_0}})$
baseband	Manchester	$Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$
passband	ASK	$Q(\sqrt{\frac{E_b}{N_0}})$
passband	PSK	$Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$
passband	FSK	$\approx Q\left(\sqrt{\frac{E_b}{N_0}}\right)$

Summary of Binary Detection



Two steps are involved in signal detection

- Step 1: reducing the received signal x(t) into a number y(T)
 - 1. find a linear filter that *maximizes* output SNR [Answer: *matched filter*]
 - 2. sample the output of filter at certain time (t = T) to get a number y(T)
- Step 2: comparing y(T) with a certain threshold and deciding which signal (i.e., $s_1(t)$ or $s_2(t)$) has been transmitted

Contents

- Noise in Communication Systems
- Optimum Detection for Binary Signals
- Probability of Error
- Summary
- Appendix

Summary of Digital Communications

1. Transmitter: Pulse-Code Modulation (PCM) [Lectures 10 & 11]



2. Channel: AWGN Channel Modeling [Lecture 12]



3. Receiver: Signal Detection [Lecture 12]



Example 1 (SNR = 0 dB): No Detection Errors



Example 1 (SNR = 0 dB): No Detection Errors

- 1. $g(t) = \cos(50\pi t) + \cos(100\pi t) \sin(33\pi t)$
- 2. samples (with $T_s = 0.009 \text{ sec}$) 2.0000 -1.5981 -1.0988 -1.3776 1.6749 1.7060
- 3. quantized samples (with 8 levels)
 1.5000 -2.0000 -1.5000 -1.5000 1.5000 1.5000
 (Note that the 8 bin centers are -2, -1.5, -1, -0.5, 0, 0.5, 1, 1.5)
- 4. bits (with 3-bit quantizer) 1 1 1, 0 0 0, 1 0 0, 1 0 0, 1 1 1, 1 1 1
- 5. transmission over AWGN channel with SNR = 0 dB
- 6. detected bits using MF

1 1 1, 0 0 0, 1 0 0, 1 0 0, 1 1 1, 1 1 1 (no errors, BER = 0)

7. detected quantized samples

1.5000 -2.0000 -1.5000 -1.5000 1.5000 1.5000

8. interpolation using sinc to get $\tilde{g}(t)$

Example 2 (SNR = -10 dB): Some Detection Errors

- 1. $g(t) = \cos(50\pi t) + \cos(100\pi t) \sin(33\pi t)$
- 2. samples (with Ts = 0.009 sec) 2.0000 -1.5981 -1.0988 -1.3776 1.6749 1.7060
- 3. quantized samples 1.5000 -2.0000 -1.5000 -1.5000 1.5000 1.5000
- 4. bits (with 3-bit quntizer)
 1 1 1, 0 0 0, 1 0 0, 1 0 0, 1 1 1, 1 1 1
- 5. transmission over AWGN channel with SNR = -10 dB
- 6. detected bits using MF 1 1 1, 0 0 0, 1 0 0, 1 0 0, 1 1 0, 1 1 0 (two errors, $\mathsf{BER}=2/18)$
- 7. detected quantized samples

 1.5000 -2.0000 -1.5000 -1.5000 -0.5000
 (two erroneous symbols, SER = 2/6)
- 8. interpolation using sinc to get $\tilde{g}(t)$

Example 2 (SNR = -10 dB): Some Detection Errors



Digital Communications vs. Analog

Main Advantage

• Digital systems are less sensitive to noise than analog

Other Advantages

- Easier to integrate different services, such as video, voice, email, etc.
- Hardware design (digital ICs are smaller and easier to make than analog ICs)
- New multiplexing techniques (TDMA, CDMA,...)
- Powerful error correcting codes and compression techniques

Main Drawback

• Digital systems typically use much higher bandwidth than analog ones **Example:** Analog telephony (3.4 kHz), Digital telephony (64 kbps)

Course Summary: What we have learned?

- 1. Fourier analysis is indispensable in communications/signal processing
 - Fourier transform table of pairs and table of properties
 - ${\ensuremath{\,\circ\,}}$ FFT implementation in ${\rm MATLAB}$ in several Labs
- 2. Filters are ubiquitous in communications/signal processing
 - Filters are used to remove unwanted signals
 - ${\ensuremath{\,\circ\,}}$ We implemented some in ${\rm MATLAB}$
- 3. Analog communication systems
 - Amplitude modulation schemes (AM, DSB, SSB, VSB)
 - Angle modulation schemes (FM, PM)
 - AM and FM radio (modulator & demodulator) circuitry
 - Superheterodyne receiver
- 4. Multiplexing techniques: FDM & TDM

Course Summary: What we have learned?

- 5. Digital communication systems
 - Sampling theorem (Nyquist rate) and interpolation
 - Quantization (linear & no-linear) and companding (μ -law & A-law)
 - Analog pulse modulation schemes (PAM, PWM, PPM)
 - Pulse code modulation (PCM) and line codes (NRZ, RZ, Manchester,...)
 - Optimum linear detection (matched filter)
 - Probability of error evaluation
- 6. Noise in communications (thermal and shot noises) and AWGN

Contents

- Noise in Communication Systems
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Matched Filter

$$SNR = \frac{|g_o(T)|^2}{\mathbb{E}[n^2(t)]} = \frac{\left| \int_{-\infty}^{\infty} H(f)G(f)e^{j2\pi fT}df \right|^2}{\frac{N_0}{2}\int_{-\infty}^{\infty} |H(f)|^2 df} \\ \leq \frac{\int_{-\infty}^{\infty} |H(f)|^2 df \times \int_{-\infty}^{\infty} |G(f)e^{j2\pi fT}|^2 df}{\frac{N_0}{2}\int_{-\infty}^{\infty} |H(f)|^2 df} \\ = \frac{2}{N_0}\int_{-\infty}^{\infty} |G(f)e^{j2\pi fT}|^2 df \\ = \frac{2}{N_0}\int_{-\infty}^{\infty} |G(f)|^2 df$$

• The inequality is due to Cauchy-Schwartz inequality¹ for $\phi_1(x) = H(f)$ and $\phi_2(x) = G(f)e^{j2\pi fT}$.

• The " \leq " will be "=" for $\phi_1(x) = \phi_2^*(x)$, i.e., for $H(f) = G^*(f)e^{-j2\pi fT}$.

 $\left| \int_{-\infty}^{\infty} \phi_1(x) \phi_2(x) dx \right|^2 \leq \int_{-\infty}^{\infty} |\phi_1(x)|^2 dx \cdot \int_{-\infty}^{\infty} |\phi_2(x)|^2 dx \text{ for any bounded complex functions } \phi_1(x) \text{ and } \phi_2(x). \text{ The equality holds if and only if } \phi_2(x) = k \phi_1^*(x).$

Lecture 12: Performance Analysis